An approach to Performance and Bottleneck Analysis

Sverre Jarp CERN openlab



Summer student lecture 2006 1 August 2006



AGENDA

- Introduction
- Path to Optimization
- Basics
- Hardware Review
- More on compilers
- Conclusion



INTRODUCTION



Initial Question

• How come we (too) often end up in the following situation?





Before we start

- This is an effort to pass in review a somewhat "systematic approach" to tuning and bottleneck analysis
 - Main focus is on understanding the "spiral to success"
- The introduction of the elements is done "topdown"
- But, it is important to understand that in real-life, this is rarely the case



PATH TO OPTIMIZATION





Step 0: Correctness

Before undertaking any tuning effort

- Excellent regression/correctness tests
 - For all critical algorithms, all important use cases

- Otherwise,

- Too many tuning efforts get left by the wayside
- Which options will work ?
 - "IPF_fp_relax"
 - "ansi_alias"
 - "ffast-math"

In an case, needed for the basic development/maintenance effort

No point in speeding up an incorrect program!

Step 1: Application design

• Regular reviews of the design (globally or partially)

- Data structures
 - Arrays; structs; data members
- Choice of algorithms
 - Accuracy, robustness, rapidity
- Design of classes
 - Domain decomposition
 - Hierarchy
 - Interrelationship
- Is there a time gap?
 - Design $\leftarrow \rightarrow$ Today's microprocessor (tomorrow's ?)
 - Did we design for low ILP, small caches, single core,....?

Step 2: Implementation aspects

Review all aspects of implementation

- Choice of language (Fortran, C, C++, Java, ...)
- Use of language features
 - Templates (STL with maps, lists, etc.)
- Precision of data (FLP)
 - Single, double, double extended
 - Intermediate calculations
 - Stored results
- Code split between .cpp and .h files
- Aggregation or decomposition ?
- Reliance on preprocessor
- Platform dependencies
 - Such as endianness
- Reliance on external libraries
 - Smartheap, Math kernel/vector libraries, etc.

1 August 2 Correct organization of source can greatly impact the application's efficiency 10



- Access to the best compiler
 - On many platforms we have a limited choice
 - IA-64/Linus or x86/MacOS: Intel or GNU (others coming?)
 - But, it is worth trying both (or all):
 - Mix and match (thanks to common ABI) ?
 - Inform the other camp when they are behind
 - Upgrade to latest versions regularly
 - Choose from hundreds of flags
- Build procedure
 - One class at a time ?
 - Archive/shared libraries ?
 - Monolithic executable or dynamic loading ?
- And (to a large extent)
 - Machine code is chosen for you

Note that x86-64 is a very healthy clean-up of the too-often extended x86 architecture



• The best hardware for the job

- Manufacturer
- Server type
 - Entry, mid-range, large SMP, NUMA, etc.
- Processor characteristics
 - Single core, Dual core, Quad core (coming)
 - Frequency, cache sizes and levels
- Further (important) factors
 - Bus speed
 - Memory speed
- Price/performance ratio

Richest choice is found inside the x86 eco-system.



In the end: Execution Results





Back to our cartoon



 As already said, first of all, we must guarantee correctness

- If we are unhappy with the performance
 - and by the way, how do we know when to be happy?
- We need to look around
 - Since the culprit can be anywhere





THE BASICS



Need a good tool set

My recommendation

- Integrated Development Environment (IDE) w/ integrated Performance Analyzer
 - Visual Studio + VTUNE (Windows)
 - Eclipse + VTUNE (Linux)
 - XCODE + Shark (MacOS)
 -

Also, other packages

- Valgrind (Linux x86, x86-64)
- Qtools (IPF)
- Pfmon, perfsuite, caliper, oprofile, TAU

Too many different tools may be counterproductive!

 Signam
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🚛 🖽 🝸 🚍 Process 🖵 Thread 🚳 Module

Performance Analyzer

VTune™

A copine Recess repo Transformed particular and a constant Transformed partice and a constant

Thread: Al

💽 View: 💿 Heavy 🔘 Tre

DISCHICS'S

Process: (49.8%) NubleApe [1705]



Price_out_impl (mcf)

• VTUNE screenshot:

Address	Line	Clockticks	Source
01:03CA	246	1426	while(arcin)
	247		{
01:03F0	248	2329	tail = arcin->tail;
01:03F2	250	323207	if(tail->time + arcin->org_cost > latest)
	251		{
01:03FC	252	1485	arcin = (arc_t *)tail->mark;
	253		continue;
	254		}
01:0410	256	657	red_cost = compute_red_cost(arc_cost, tail, head_potential);
01:0424	258	2233	if(red_cost < 0)
	259		{
01:04F6	260	296	if(new_arcs < MAX_NEW_ARCS)
	261		{
	262		insert_new_arc(arcnew, new_arcs, tail, head,
01:0530	263	39	arc_cost, red_cost);"
01:0558	264	16	new_arcs++;
	265		}
01:04FE	266	301	else if((cost_t)arcnew[0].flow > red_cost)
	267		replace_weaker_arc(arcnew, tail, head,
01:050B	268	65	arc_cost, red_cost);
	269		}
01:042C	271	1664	arcin = (arc_t *)tail->mark;
	272		}



- The language spoken by the processor is
 MACHINE CODE !!
- To understand it, we need what I call "Assembler awareness":
 - Looking into compiler-generated code, there may be a need to:
 - Modify (repeatedly) the HLL code (or compiler options) and inspect the result
 - When available, add inline assembly or intrinsics for localized impact
 - Today, we are not dealing with the case of writing Assembly code
 - But the issues are the same



Machine code

It may be necessary to readirectly

```
Bool_t TGeoCone::Contains(Double_t *point) const
{
    // test if point is inside this cone
    if (TMath::Abs(point[2]) > fDz) return kFALSE;
```

```
Double_t r2 = point[0]*point[0] + point[1]*point[1];
Double_t rl = 0.5^{(fRmin2^{(point[2] + fDz) + fRmin1^{(point[2] + fDz) + fRmax)}
Double_t rh = 0.5^{(fRmax2^{(point[2] + fDz) + fRmax)}
if ((r2<rl*rl) || (r2>rh*rh)) return kFALSE;
return kTRUE;
```

ZNK8TGeoCone8ContainsEPd: [.LFB1785:] .proloque .body .mmi adds r14 = 16, r33adds r15 = 16, r32adds r16 = 32, r32. mmi adds r17 = 24, r32adds r18 = 40, r32adds r32 = 8, r32 ;; .mmi ldfd f11 = [r14]1dfd f15 = [r32]mov r8 = r0;; .mfb $fcmp.ge \ p6, \ p7 = f11, \ f0$.mfi mov f6 = f11;; .mmf (p7) fneg f6 = f11 ;; .mmf fcmp.gt p6, p7 = f6, f15;; .bbb (p6) br.ret.dptk.many rp

(snip)

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Profiling with q-tools

• Test40: Physics simulation job

Command: ./test40icc8002fz Flat profile of CPU_CYCLES in test40icc8002fz-pid24686-cpu0.hist#0: Each histogram sample counts as 1.00034m seconds % time self cumul calls self/call tot/call name 14.28 7.23 7.23 1.72M 4.21u 4.90u G4VoxelNavigation::LevelLocate(G4NavigationHistory&, G4VPhysicalVolume const*, int, Hep3Vector const&, Hep3Vector const*, bool, Hep3Vector&) 10.08 35.4M 80.5n 80.5n RanecuEngine::flat() 5.63 2.85 476n 2.94u G4Navigator::LocateGlobalPointAndSetup(Hep3Vector const&, Hep3Vector const*, 3.57 1.81 11.89 3.80M bool,...) 3.46 1.75 13.64 5.12M 343n 2.54u G4SteppingManager::DefinePhysicalStepLength() 2.30 14.80 898k 1.30u 2.84u G4VEnergyLoss::GetLossWithFluct(G4DynamicParticle const*, G4Material*, double,) 1.16 2.29 41.2n G4Tubs::Inside(Hep3Vector const&) const 1.16 15.96 28.1M 41.2n 2.27 1.15 7.50M 153n 2.27u G4SteppingManager::InvokePSDIP(unsigned long) 17.11 2.23 1.13 18.24 5.17M 219n 7.53u G4SteppingManager::Stepping() 2.17 1.10 19.34 4.92M 223n 2.03u G4Transportation::PostStepDoIt(G4Track const&, G4Step const&) 2.12 1.08 15.8M 67.9n 67.9n G4PhysicsLogVector::FindBinLocation(double) const 20.41 1.93 21.39 5.23M 186n 901n G4Transportation::AlongStepGetPhysicalInteractionLength(G4Track const&, double, 0.98 double, double&, G4GPILSelection*) 1.90 22.35 1.11M 864n 864n G4MuPairProduction::ComputeDDMicroscopicCrossSection(G4ParticleDefinition const*,... 0.96) 1.80 0.91 23.26 1.01M 897n 1.50u G4MultipleScattering::PostStepDoIt(G4Track const&, G4Step const&) 24.16 47.3n 47.3n G4Track::GetVelocity() const 1.78 0.90 19.1M 1.62 0.82 24.98 4.48M 182n 720n G4Navigator::ComputeStep(Hep3Vector const&, Hep3Vector const&, double, double&) 863n G4MultipleScattering::GetContinuousStepLimit(G4Track const&, double, double, double&) 1.51 0.77 25.75 1.89M 405n 1.41 0.71 26.46 4.62M 155n 503n G4ReplicaNavigation::ComputeStep(Hep3Vector const&, Hep3Vector const&, Hep3Vector const&, Hep3Vector const&, double, double&, G4NavigationHistory&, bool&, Hep3Vector&, bool&, bool&, G4VPhysicalVolume**, int&) 27.10 3.27M 196n 196n G4Tubs::DistanceToOut(Hep3Vector const&, Hep3Vector const&, bool, bool*, 1.27 0.64 Hep3Vector*) const 1.08 0.55 27.65 550k 999n 2.28u G4eBremsstrahlung::PostStepDoIt(G4Track const&, G4Step const&) 1.07 0.54 28.20 5.01M 109n 170n G4Transportation::AlongStepDoIt(G4Track const&, G4Step const&) 70.7n 70.7n int malloc 0.98 0.50 28.69 7.00M 102n 0.96 0.49 29.18 4.80M 3.65u G4SteppingManager::InvokePostStepDoItProcs() 0.93 29.65 18.2u 18.6u G4MultipleScattering::ComputeTransportCrossSection(G4ParticleDefinition const&,) 0.47 25.8k 93.8n 146n Em2SteppingAction::UserSteppingAction(G4Step const*) 0.92 0.47 30.12 4.97M 0.91 30.58 4.95M 93.4n 1.04u G4SteppingManager::InvokeAlongStepDoItProcs() 0.46 0.91 230n G4VDiscreteProcess::PostStepGetPhysicalInteractionLength(G4Track const&, double, 0.46 31.04 9.64M 47.5n



HARDWARE REVIEW



CPU performance vector







Memory Hierarchy





Cache lines

- Madison L3 cache lines are 128B (16 * double)
 - Minimum amount of data transferred between cache and memory.
 - Imagine what happens if your stride is 16 (or more)!

Programming the memory hierarchy is an art in itself.



Back to Compilers





1 August 2006

	TestKalman [nx,ny] : kalman_win7.1								
_	2	3	4	5	6	7	8	9	10
2	0.41	0.55	0.80	1.26	2.16	3.91	5.70	7.43	9.66
	0.40	0.56	0.79	1.31	2.36	3.84	5.14	6.97	8.90
	1.01	1.29	1.65	2.18	3.16	7.13	8.85	11.24	13.66
3	0.52	0.70	0.98	1.49	2.43	4.35	6.07	8.23	10.09
	0.51	0.70	0.98	1.52	2.62	4.15	5.46	7.55	9.17
	1.24	1.50	1.97	2.61	3.68	7.82	9.53	12.26	14.95
4	0.63	0.85	1.24	1.73	2.79	4.77	6.65	8.64	10.86
	0.62	0.85	1.17	1.77	2.86	4.46	5.93	7.93	10.01
	1.50	1.88	2.19	3.09	4.31	8.53	10.56	13.77	16.58
5	0.78	1.04	1.41	2.11	3.12	5.12	7.17	9.64	11.45
	0.83	1.09	1.45	2.10	3.22	4.90	6.53	8.70	10.56
	1.81	2.24	2.91	3.49	5.02	9.40	11.56	14.88	17.61
6	0.85	1.16	1.68	2.28	3.50	5.57	8.12	9.94	12.50
	0.98	1.29	1.72	2.49	3.72	5.44	7.07	9.22	11.42
	2.13	2.65	3.40	4.37	5.49	10.36	12.53	16.09	19.24
7	1.04	1.50	2.01	2.79	4.03	6.24	8.48	10.76	13.30
	1.10	1.48	1.99	2.80	4.15	5.89	7.64	9.80	11.96
	2.44	3.09	3.95	4.95	6.47	10.88	13.44	17.59	20.76
8	1.22	1.69	2.30	3.18	4.59	6.89	9.24	11.67	14.35
	1.26	1.71	2.30	3.16	4.57	6.47	8.69	10.78	13.03
	2.81	3.57	4.48	5.57	7.28	12.02	14.23	18.69	22.77
N1,N2 <= 6 36.51 37.96 66.81 N1,N2 > 6 261.15 242.13 421.54 All N1,N2 297. 480. 488.					297.67 280.08 488.35				
SM	SMatrix_Sym SMatrix TMatrix SMatrix_Sym better than TMatrix								

TestKalman [nx,ny] : kalman_slc3_gcc323									
	2	3	4	5	6	7	8	9	10
2	0.30	0.38	0.64	0.88	1.64	5.84	7.81	10.41	14.36
	0.86	1.00	1.39	1.69	2.96	5.22	6.06	7.41	9.03
2	0.37	0.56	0.72	1.10	1.84	6.24	8.19	10.97	14.37
3	1.01	1.16	1.59	1.96	3.40	5.48	6.64	7.61	9.86
	0.47	0.63	0.89	1.39	2.16	6.71	9.04	11.71	15.07
4	0.61	0.76	1.03	1.48	2.27	6.43 6.14	8.92 6.95	11.37 8.85	14.69
	0.60	0.85	1.19	1.71	2.58	7.03	9.52	12.41	15.74
5	0.78	1.03	1.28	1.80	2.70	6.79	9.18	12.03	16.00
	1.28	1.55	2.35	2.69	4.44	6.60	8.34	9.44	12.16
6	0.77	1.26	1.49	2.13	3.06	7.81	10.19 9.90	12.98	17.56
Ŭ	1.59	2.09	2.42	3.58	4.61	7.55	8.09	10.36	11.58
_	0.96	1.33	1.77	2.46	3.47	8.24	10.56	13.08	18.03
7	1.25	1.49	1.99	2.57	3.62 5.48	8.08 7.90	10.13 9.50	12.72 10.62	16.96 14.14
	1.14	1.68	2.15	2.95	4.07	8,99	11.47	14.48	19,15
8	1.48	1.79	2.33	3.14	4.27	8.79	11.37	14.36	18.07
	2.05	2.81	3.02	4.67	5.59	8.69	9.34	12.29	13.71
			29.45		_	340.08			369.53
N1,N2 <= 6		2 <= 6 32.23 N1,N2 > 6		2 > 6	334.29 All N1,N2			366.52	
54.07 283.99 338.05						338.05			
SMatrix_Sym SMatrix TMatrix SMatrix_Sym better than TMatrix									

TestKalman [nx,ny] : kalman_solaris.5.9

	2	3	4	5	6	7	8	9	10
	2.29	4.53	7.49	13.36	27.92	30.68	42.12	56.27	74.79
2	1.39	2.41	3.52	5.57	9.88	20.13	29.90 42.52	41.89	56.08 61 38
	0.00	2.35	0.04	47.00	0.10	07.50	54.07	01.00	01.00
3	2.05	3.43	9.88 5.09	7.79	12.81	24.26	35.10	50.11	66.09
•	2.83	3.49	4.74	6.23	9.61	36.25	44.61	53.69	63.78
	4.70	8.39	13.02	21.30	38.09	44.86	62.55	83.96	108.25
4	2.92	4.82	7.16	10.45	16.27	28.23	43.15	60.94	79.72
	3.50	4.41	5.68	7.46	11.42	38.35	47.23	56.35	67.32
_	6.45	11.09	16.75	26.01	45.35	52.92	73.96	100.57	127.69
5	3.84	6.42	9.42	13.43	20.22	32.57	50.84	72.06	94.24
	3.87	5.10	6.78	8.88	12.96	40.86	49.51	59.60	70.80
-	8.77	14.55	21.27	32.27	51.35	63.23	87.57	118.44	152.52
6	5.36	8.67	12.45	17.49	25.37	39.55	60.54	84.29	112.31
	4.58	6.12	8.33	10.55	14.90	43.39	52.76	63.18	75.01
_	12.58	20.21	29.16	42.12	64.82	78.81	107.49	142.03	183.05
7	6.85	10.88	15.45	21.34	29.96	44.96	68.27	96.24	128.52
	5.27	7.13	9.41	12.30	17.47	45.91	55.53	66.99	79.38
-	17.68	28.33	40.40	57.23	84.57	103.45	139.67	184.39	232.32
8	10.79	17.19	24.55	33.55	46.08	64.54	95.46	132.12	170.91
	6.08	8.30	10.98	14.20	19.58	48.60	58.98	/0.5/	83.94
			445.38			3095.72			3541.10
	N1,N	2 <= 6	218.23	N1,N	2 > 6	2099.65	All N	1,N2	2317.89
	,		164.08	,		1673.16			1837.24
SN	SMatrix_Sym SMatrix TMatrix SMatrix_Sym better than TMatrix								

Lecture - Sourtesy: René Brun/CERN



"Low-hanging fruit"

• Typically one starts with a given compiler, and moves to:

Summer Student Lecture - SJ

- More aggressive compiler options
 - For instance:
 - -O2 → -O3,-funroll-loops, -ffast-math (g
 - -O2 → -O3, -ipo (icc)

Some options can compromise accuracy or correctness

- More recent compiler versions
 - g++ version 3 \rightarrow g++ version 4
 - icc version $8 \rightarrow$ icc version 9
- Different compilers
 - GNU \rightarrow Intel (or reverse?)

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May be a burden because of potential source code issues



Interprocedural optimization

- Let the compiler worry about interprocedural relationship
 - "icc --ipo"
- Valid also when building libraries
 - Archive
 - Shared
- Cons:
 - Can lead to code bloat
 - Longer compile times

Probably most useful when combined with heavy optimization for "production" binaries or libraries!





Feedback Optimization

- Many compilers allow further optimization through training runs
 - Compile once (to instrument binary)
 - g++ -fprofile-generate
 - icc -prof_gen
 - Run one (or several test cases)
 - ./test40 < test40.in (will run slowly)
 - Recompile w/feedback
 - g++ -fprofile-use
 - icc -prof_use (best results when combined with -O3,-ipo

With icc 9.0 we get ~20% on root stress tests on Itanium, but only ~5% on x86-64



CONCLUSION



Conclusions

- Understand which parts of the "spiral" you control
- Understand the platform hardware
- Equip yourself with good tools
 - Get access to hw performance counters
 - Exploit the power of performance tools
- Check how key algorithms map on to your hardware platform
 - Are you at 5% or 95% efficiency?
 - Where do you want to be?
- Cycle around the spiral frequently
 - It is hard to get to "peak" performance (and stay there!)



QUESTIONS?



Backup



In comes the PMU (Performance Monitoring Unit)

Quickly summarized: 4 counters (12 on Montecito) ~200 monitored events Some very advanced features!



Itanium-2 cache hierarchy



37



Geant 4 – Test40

• Overall counters in 10^9

Counter	Counts
IA64_INST_RETIRED	108.43
NOPS_RETIRED	39.73
CPU_CYCLES (CC)	74.98
BACK_END_BUBBLE_ALL	43.10

Useful instructions (UI)	68.70
Non-stalled cycles (NSC)	31.88

UI/CC	~ 1
UI/NSC	~ 2



Geant 4 – Test40

• Stall counters

Counter	Counts (10^9)	
BACK_END_BUBBLE_ALL	43.10	
BE_EXE_BUBBLE_ALL	27.96	
BACK_END_BUBBLE_FE	7.38	
BE_L1D_FPU_BUBBLE_ALL	3.86	
BE_FLUSH_BUBBLE_ALL	2.72	
BE_RSE_BUBBLE_ALL	1.38	

Why so many EXE bubbles?

Keep drilling down!

65

%



Geant 4 – Test40

• EXE stall counters

Counter	Counts
BE_EXE_BUBBLE_ALL	27.96
BE_BUBBLE_GRALL	10.29
BE_BUBBLE_GRGR	~zero
BE_EXE_BUBBLE_FRALL	17.21

Counter	Counts
L2_REFERENCES	19.39
L2_DATA_REFERENCES_L2_ALL	13.16

ldf,ldf
bubble
fma
bubble
bubble
bubble
stf



Test40 - Cache counters





Software Pipelining



}

Mersenne Twister

```
Double t TRandom3::Rndm(Int t){
    UInt t y;
    const Int t kM = 397; const Int t kN = 624; const UInt t kTemperingMaskB = 0x9
    const UInt_t kTemperingMaskC = 0xefc60000; const UInt t kUpperMask =
                                                                                  0 \times 800
    const UInt t kLowerMask = 0x7fffffff; const UInt t kMatrixA =
                                                                                  0x990
    if (fCount624 >= kN) {
       register Int t i;
      for (i=0; i < kN-kM; i++) { /* THE LOOPS */</pre>
        y = (fMt[i] & kUpperMask) | (fMt[i+1] & kLowerMask);
        fMt[i] = fMt[i+kM] \wedge (y >> 1) \wedge ((y \& 0x1) ? kMatrixA : 0x0);
      }
      for ( ; i < kN-1 ; i++) {
        y = (fMt[i] & kUpperMask) | (fMt[i+1] & kLowerMask);
        fMt[i] = fMt[i+kM-kN] ^ (y >> 1) ^ ((y \& 0x1) ? kMatrixA : 0x0);
      }
      y = (fMt[kN-1] & kUpperMask) | (fMt[0] & kLowerMask);
      fMt[kN-1] = fMt[kM-1] ^{(y >> 1)} ^{((y \& 0x1))} RMatrixA : 0x0);
      fCount624 = 0;
    }
    y = fMt[fCount624++]; /*THE STRAIGHT-LINE PART*/
    y^{+} (y >> 11); y^{+} ((y << 7)) \& kTemperingMaskB);
    y ^= ((y << 15) & kTemperingMaskC ); y ^= (y >> 18);
    if (y) return ( (Double t) y * 2.3283064365386963e-10); // * Power(2,-32)
    return Rndm();
```



The "MT" loop is full

• Highly optimized

- Here depicted in 3 Itanium cycles
 - But similarly dense on other platforms

0	Load	Test Bit	XOR	Load	Add	No-op
1	AND	AND	Shift	Add	Load	Move
2	Store	OR	XOR	Add	Add	Branch

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The sequential part is not!

	0	Add	Mov long	No-op	No-ор	No-op	No-op
	1	Load	Mov long	Mov long	No-ор	No-op	No-op
	2	Shift,11	Set float	No-op	No-ор	No-op	No-op
	3	XOR	Move	No-op	No-ор	No-op	No-ор
	4	Shift,7	No-op	No-op	No-ор	No-op	No-ор
	5	AND	No-op	No-op	No-ор	No-op	No-ор
	6	XOR	No-op	No-op	No-ор	No-op	No-op
	7	SHL,15	No-op	No-op	No-ор	No-op	No-op
	8	AND	No-op	No-op	No-ор	No-op	No-op
	9	XOR	No-op	No-op	No-ор	No-op	No-op
	10	SHL,18	No-op	No-op	No-ор	No-op	No-op
	11	XOR	No-op	No-op	No-ор	No-op	No-op
	12	Set float	Compare	Branch	No-ор	No-op	No-op
	13	Bubble (no work disp $y = fMt[fCount624++]; /*THE STRAIGHT-LINE PART*/$					
14 Bubble (no work disp $y^{+} = (y >> 11); y^{+} = ((y << 7)) \& k$) & kTemper	ingMaskB);
	15 Bubble (no work disp. $y^{+} = ((y << 15) \& kTemperingMaskC); y^{+} =$						(y >> 18);
16 Bubble (no work disp if (y) return ((Double_t) y * 2.3283064						2.3283064365	386963e-10);
	17	Bubble (no work dispatched, because of FP latency)					
	18	Mult FP	No-op	No-op	No-op	No-op	No-op
	19	Bubble (no work dispatched, because of FP latency)					
	20	Bubble (no work dispatched, because of FP latency)					
	21 Bubble (no work dispatched, because of FP latency)						
	22	Mult FP	Branch	No-op	No-op	No-op	No-op
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